

Hutch++

Optimal Stochastic Trace Estimation

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Trace Estimation

- ⊙ Goal: Estimate trace of $d \times d$ matrix \mathbf{A} :

$$\text{tr}(\mathbf{A}) = \sum_{i=1}^d \mathbf{A}_{ii} = \sum_{i=1}^d \lambda_i$$

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- Instead, \mathbf{B} is in memory and $\mathbf{A} = f(\mathbf{B})$:

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- Goal: Estimate $\text{tr}(\mathbf{A})$ by computing $\mathbf{A}\mathbf{x}_1, \dots, \mathbf{A}\mathbf{x}_k$

Implicit Matrix Trace Estimation: Estimate $\text{tr}(\mathbf{A})$ with as few Matrix-Vector products $\mathbf{A}\mathbf{x}_1, \dots, \mathbf{A}\mathbf{x}_m$ as possible.

$$(1 - \varepsilon) \text{tr}(\mathbf{A}) \leq \tilde{\text{tr}}(\mathbf{A}) \leq (1 + \varepsilon) \text{tr}(\mathbf{A}) \quad \text{w.p. } 1 - \delta$$

For PSD matrix trace estimation,

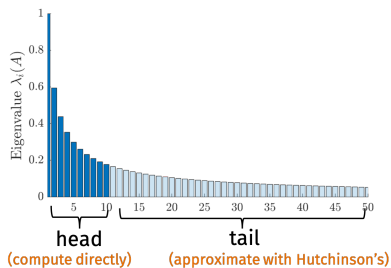
- ⊙ Hutch++ algorithm, which uses $\tilde{O}(\frac{1}{\epsilon})$ matrix-vector products.
 - Improves prior rate of $\tilde{O}(\frac{1}{\epsilon^2})$
 - Empirically works well
 - Matching $\Omega(\frac{1}{\epsilon})$ Lower Bound

Only 5 lines of code:

```
1 function T = hutchplusplus(A, m)
2     S = 2*randi(2,size(A,1),m/3);
3     G = 2*randi(2,size(A,1),m/3);
4     [Q,~] = qr(A*S,0);
5     G = G - Q*(Q'*G);
6     T = trace(Q'*A*Q) + 1/size(G,2)*trace(G'*A*G);
7 end
```

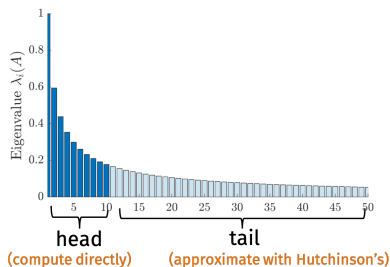
¹ \tilde{O} notation only hide logarithmic dependence on the failure probability.

Hutch++ Intuition



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1. Find a good rank- k approximation $\tilde{\mathbf{A}}_k$
2. Compute $\tilde{T} \approx \text{tr}(\mathbf{A} - \tilde{\mathbf{A}}_k)$ with m steps of Hutchinson's
3. Return $\text{Hutch++}(\mathbf{A}) = \text{tr}(\tilde{\mathbf{A}}_k) + \tilde{T}$

THANK
YOU

Code available at
github.com/RaphaelArkadyMeyerNYU/hutchplusplus